

Elimination Theory

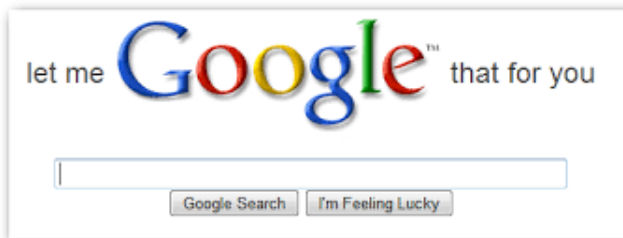
Carlos D'Andrea

Fourth EACA International School on Computer Algebra and
its Applications



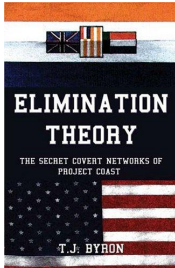
What is Elimination Theory?

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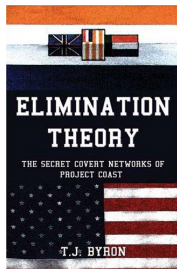


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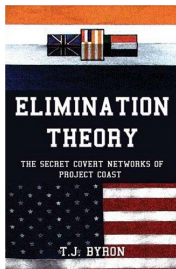


Elimination Theory



$$\operatorname{Res}(f(x), g(x)) = \begin{vmatrix} a_n & a_{n-1} & \cdots & a_1 & a_0 & 0 & \cdots & 0 \\ 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 \\ b_n & b_{n-1} & \cdots & b_1 & b_0 & 0 & \cdots & 0 \\ 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 \end{vmatrix}.$$

Elimination Theory



$\text{Res}(f(x), g(x)) =$

$$\begin{vmatrix} a_n & a_{n-1} & \cdots & a_1 & a_0 & 0 & \cdots & 0 \\ 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 \\ b_n & b_{n-1} & \cdots & b_1 & b_0 & 0 & \cdots & 0 \\ 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 \end{vmatrix}$$

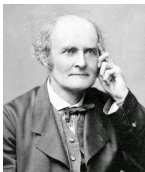


A bit of history

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In the problem of **elimination**, one seeks the relationship that must exist between the coefficients of a function or system of functions in order that some particular circumstance (or singularity) can occur.

Cayley, 1864 Nouvelles recherches sur l'élimination et la théorie des courbes



A bit of history

A system of arbitrarily many algebraic equations for $z^0, z', z'', \dots, z^{(n-1)}$, in which the coefficients belong to the rationality domain $(R, R', R'', \dots)$, define the algebraic relations between z and R , whose knowledge and representation are the purpose of the **theory of elimination**.

Kronecker, 1882 Grundzüge einer arithmetischen Theorie der algebraischen Größen



The baby example in elimination

The baby example in elimination

Find “the condition” on
 $a_{10}, a_{11}, a_{20}, a_{21}$ so that the system

$$\begin{cases} a_{10}x_0 + a_{11}x_1 = 0 \\ a_{20}x_0 + a_{21}x_1 = 0 \end{cases}$$

has a solution different from $(0, 0)$

“Elimination”

$$a_{10}x_0 + a_{11}x_1, a_{20}x_0 + a_{21}x_1$$

$$\in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}, x_0, x_1]$$

“Elimination”

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$$\in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}, x_0, x_1]$$

↓

$$a_{10}a_{21} - a_{20}a_{11} \in \mathbb{K}[a_{10}, a_{11}, a_{20}, a_{21}]$$

More general

Find “the condition” for the system

$$\left\{ \begin{array}{lcl} a_{10}x_0 + a_{11}x_1 + \dots + a_{1n}x_n & = & 0 \\ a_{20}x_0 + a_{21}x_1 + \dots + a_{2n}x_n & = & 0 \\ \vdots & & \vdots \\ a_{(n+1)0}x_0 + a_{(n+1)1}x_1 + \dots + a_{(n+1)n}x_n & = & 0 \end{array} \right.$$

More general

Find “the condition” for the system

$$\left\{ \begin{array}{lcl} a_{10}x_0 + a_{11}x_1 + \dots + a_{1n}x_n & = & 0 \\ a_{20}x_0 + a_{21}x_1 + \dots + a_{2n}x_n & = & 0 \\ \vdots & & \vdots \\ a_{(n+1)0}x_0 + a_{(n+1)1}x_1 + \dots + a_{(n+1)n}x_n & = & 0 \end{array} \right.$$

to have a solution different from
 $(0, 0, \dots, 0)$

Another more general

Let $d_1, d_2 \in \mathbb{N}$. Find “the condition”
for the system of polynomials

$$\begin{cases} a_{10}x_0^{d_1} + a_{11}x_0^{d_1-1}x_1 + \dots = 0 \\ a_{20}x_0^{d_2} + a_{21}x_0^{d_2-1}x_1 + \dots = 0 \end{cases}$$

to have a solution different from
 $(0, 0)$

More more more general...

Let $n \in \mathbb{N}$, and $d_1, \dots, d_{n+1} \in \mathbb{N}$, find the condition
for

$$\left\{ \begin{array}{lcl} \sum_{\alpha_0 + \dots + \alpha_n = d_1} a_{1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \\ \sum_{\alpha_0 + \dots + \alpha_n = d_2} a_{2, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \\ \vdots & & \vdots \\ \sum_{\alpha_0 + \dots + \alpha_n = d_{n+1}} a_{n+1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} & = & 0 \end{array} \right.$$

to have a solution different from $(0, 0, \dots, 0)$

Elimination: The general problem

For $\mathbf{a} = (a_1, \dots, a_N)$, $k, n \in \mathbb{N}$ let
 $f_1(\mathbf{a}, x_1, \dots, x_n), \dots, f_k(\mathbf{a}, x_1, \dots, x_n) \in$
 $\mathbb{K}[\mathbf{a}, x_1, \dots, x_n]$. Find conditions on $\mathbf{a}$ such that

$$\left\{ \begin{array}{lll} f_1(\mathbf{a}, x_1, \dots, x_n) & = & 0 \\ f_2(\mathbf{a}, x_1, \dots, x_n) & = & 0 \\ \vdots & \vdots & \vdots \\ f_k(\mathbf{a}, x_1, \dots, x_n) & = & 0 \end{array} \right.$$

has a solution

Solution?

- Depends on the ground field

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- There is not necessarily a “closed” condition

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- Depends on the ground field
- There is not necessarily a “closed” condition
- Algorithms??

Easy example

$$\left\{ \begin{array}{lcl} a_{11}x_1 + \dots + a_{1n}x_n & = & 0 \\ a_{21}x_1 + \dots + a_{2n}x_n & = & 0 \\ \vdots & & \vdots \\ a_{kn}x_1 + \dots + a_{kn}x_n & = & 0 \end{array} \right.$$

with $k \geq n$

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with $k \geq n$

Conditions: all maximal minors of $(a_{ij})_{1 \leq i \leq k, 1 \leq j \leq n}$
equal to zero

Another “easy” example

$$k = n = 1,$$

$$a_0 + a_1x_1 + a_2x_1^2 + \dots + a_dx_1^d = 0$$

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Conditions?

Geometry

$$V = \{(\mathbf{a}, x_1, \dots, x_n) : f_1(\mathbf{a}, x_1, \dots, x_n) = 0, \dots, f_k(\mathbf{a}, x_1, \dots, x_n) = 0\}$$

Geometry

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$$V \subset \mathbb{K}^N \times \mathbb{K}^n$$

$$\pi_1|_V \downarrow \qquad \downarrow \pi_1$$

$$\pi_1(V) \subset \mathbb{K}^N$$

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The set of conditions is $\pi_1(V)$ is not necessarily described by zeroes of polynomials

Projective Elimination

(homogeneous polynomials in an algebraically closed field)

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$$\pi_1(V) = \{p_1(\mathbf{a}) = 0, \dots, p_\ell(\mathbf{a}) = 0\}$$

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Mertens 1889



Mertens 1889



Set $s := \max \deg(f_i)$, and $G_0, \dots, G_n$ expressions of the form

$G_j(\mathbf{a}, \mathbf{x}) := \sum u_{\alpha,i} \mathbf{x}^\alpha f_i(\mathbf{a}, \mathbf{x})$ of degree s in $\mathbf{x}$ with new parameters $u_{\alpha,i}$.

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Then, “the Kronecker $\mathbf{u}$ -Resultant”

$\text{Res}_{\mathbf{x}}(G_0, \dots, G_n) = \sum_{\beta} P_{\beta}(\mathbf{a}) \mathbf{u}^{\beta}$ gives the conditions $\{P_{\beta}(\mathbf{a})\}$

van der Waerden 1926



van der Waerden 1926



For arbitrary s , we use the “Macaulay matrix”

$M_s(\mathbf{a})$ such that

$$M_s(\mathbf{a}) \begin{bmatrix} \vdots \\ \mathbf{x}^\beta \\ \vdots \end{bmatrix} = \begin{bmatrix} \vdots \\ \mathbf{x}^\alpha f_i \\ \vdots \end{bmatrix}, \quad |\alpha| = s$$

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For $s \gg 0$, all maximal minors of $M_s(\mathbf{a})$ vanish iff there is a common solution

Algebra & Geometry

(van der Waerden, 1926)

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$$\begin{aligned} I(\pi_1(V)) \\ = \\ (\langle f_1, \dots, f_k \rangle : \langle x_0, \dots, x_n \rangle^\infty) \end{aligned}$$

Algebra & Geometry

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$$\begin{aligned} & I(\pi_1(V)) \\ &= \\ & (\langle f_1, \dots, f_k \rangle : \langle x_0, \dots, x_n \rangle^\infty) \cap \mathbb{K}[\mathbf{a}] \end{aligned}$$

The “elimination ideal”

But what if...



But what if...



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- the polynomials are not homogeneous?
- $\mathbb{K}$ is not algebraically closed?

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- . . .

Homogeneous vs non homogeneous

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Analogy with Linear Algebra...

$$\begin{array}{c} \left[\begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 0 & 0 & -1 & -1 \\ 0 & 1 & 1 & 2 \end{array} \right] \\ \downarrow \\ \left[\begin{array}{ccc|c} 1 & 1 & 1 & 3 \\ 0 & 1 & 1 & 2 \\ 0 & 0 & -1 & -1 \end{array} \right] \end{array}$$

Generalization of the determinant?

$$V = \{(\mathbf{a}, x_0, x_1, \dots, x_n) : f_1(\mathbf{a}, x_0, x_1, \dots, x_n) = 0, \dots, f_{n+1}(\mathbf{a}, x_0, x_1, \dots, x_n) = 0\}$$

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One hopes!

Example 1

$$\left\{ \begin{array}{l} a_{00}x_0 + a_{01}x_1 + \dots + a_{0n}x_n = 0 \\ a_{10}x_0 + a_{11}x_1 + \dots + a_{1n}x_n = 0 \\ \qquad \qquad \qquad \vdots \qquad \qquad \qquad \vdots \qquad \vdots \\ a_{n0}x_0 + a_{n1}x_1 + \dots + a_{nn}x_n = 0 \end{array} \right.$$

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$$p_1(a) = \det(a_{ij})$$

Example 2

$$\begin{cases} f_1 = a_{10}x_0^{d_1} + a_{11}x_0^{d_1-1}x_1 + \dots + a_{1d_1}x_1^{d_1} \\ f_2 = a_{20}x_0^{d_2} + a_{21}x_0^{d_2-1}x_1 + \dots + a_{2d_2}x_1^{d_2} \end{cases}$$

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$$\text{Res}(f_1, f_2) = \det \begin{pmatrix} a_{10} & a_{11} & \dots & a_{1d_1} & 0 & \dots & 0 \\ 0 & a_{10} & \dots & a_{1d_1-1} & a_{1d_1} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{10} & \dots & \dots & a_{1d_1} \\ a_{20} & a_{21} & \dots & a_{2d_2} & 0 & \dots & 0 \\ 0 & a_{20} & \dots & a_{2d_2-1} & a_{2d_2} & \dots & 0 \\ \vdots & \vdots & \ddots & \dots & \dots & \ddots & \vdots \\ 0 & 0 & \dots & a_{20} & \dots & \dots & a_{2d_2} \end{pmatrix}$$

Example 3

$$\left\{ \begin{array}{l} f_1 = \sum_{\alpha_0 + \dots + \alpha_n = d_1} a_{1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ f_2 = \sum_{\alpha_0 + \dots + \alpha_n = d_2} a_{2, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \\ \vdots \\ f_{n+1} = \sum_{\alpha_0 + \dots + \alpha_n = d_{n+1}} a_{n+1, \alpha_0, \dots, \alpha_n} x_0^{\alpha_0} \dots x_n^{\alpha_n} \end{array} \right.$$

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$$\text{Res}_{d_1, \dots, d_n}(f_1, f_2, \dots, f_{n+1})$$

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$\text{Res}_{d_1, \dots, d_n}(f_1, f_2, \dots, f_{n+1})$
Macaulay/dense/classical resultant

No-example

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$$\left\{ \begin{array}{l} a_{00}x_0^4 + a_{01}x_0^2x_1x_2 + a_{02}x_1^2x_2^4 = 0 \\ a_{10}x_0^4 + a_{11}x_0^2x_1x_2 + a_{12}x_1^2x_2^4 = 0 \\ a_{20}x_0^4 + a_{21}x_0^2x_1x_2 + a_{22}x_1^2x_2^4 = 0 \end{array} \right.$$

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$$p_1(\mathbf{a}) \neq \det(a_{ij})$$

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$$p_1(\mathbf{a}) \neq \det(a_{ij})$$

$$p_1(\mathbf{a}) = a_{00}^2a_{12}^2 - a_{00}a_{01}a_{11}a_{12} - 2a_{00}a_{02}a_{10}a_{12} + a_{00}a_{02}a_{11}^2 + a_{01}^2a_{10}a_{12} - a_{01}a_{02}a_{10}a_{11} + a_{02}^2a_{10}^2$$

No-example

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$$p_2(\mathbf{a}) = a_{00}^2a_{22}^2 - a_{00}a_{01}a_{21}a_{22} - 2a_{00}a_{02}a_{20}a_{22} + a_{00}a_{02}a_{21}^2 + a_{01}^2a_{20}a_{22} - a_{01}a_{02}a_{20}a_{21} + a_{02}^2a_{20}^2$$

No-example

$$\begin{cases} a_{00}x_0^4 + a_{01}x_0^2x_1x_2 + a_{02}x_1^2x_2^4 = 0 \\ a_{10}x_0^4 + a_{11}x_0^2x_1x_2 + a_{12}x_1^2x_2^4 = 0 \\ a_{20}x_0^4 + a_{21}x_0^2x_1x_2 + a_{22}x_1^2x_2^4 = 0 \end{cases}$$

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$$p_1(\mathbf{a}) = a_{00}^2a_{12}^2 - a_{00}a_{01}a_{11}a_{12} - 2a_{00}a_{02}a_{10}a_{12} + a_{00}a_{02}a_{11}^2 + a_{01}^2a_{10}a_{12} - a_{01}a_{02}a_{10}a_{11} + a_{02}^2a_{10}^2$$

$$p_2(\mathbf{a}) = a_{00}^2a_{22}^2 - a_{00}a_{01}a_{21}a_{22} - 2a_{00}a_{02}a_{20}a_{22} + a_{00}a_{02}a_{21}^2 + a_{01}^2a_{20}a_{22} - a_{01}a_{02}a_{20}a_{21} + a_{02}^2a_{20}^2$$

$$p_3(\mathbf{a}) = a_{20}^2a_{12}^2 - a_{20}a_{21}a_{11}a_{12} - 2a_{20}a_{22}a_{10}a_{12} + a_{20}a_{22}a_{11}^2 + a_{21}^2a_{10}a_{12} - a_{21}a_{22}a_{10}a_{11} + a_{22}^2a_{10}^2$$

Cautionary Example

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“Triangulating” a system

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“Triangulating” a system

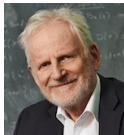
Two ways:

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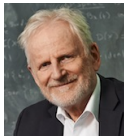
Both use some kind of **division** of
polynomials

$$\begin{array}{rcl} a_1: & & y \\ a_2: & & -1 \\ xy+1 & & \\ y+1 & & \sqrt{xy^2+1} \\ & & \hline & & xy^2+y \\ & & -y+1 \\ & & \hline & & -y-1 \\ & & \hline \end{array}$$

Gröbner Bases

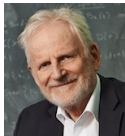


Gröbner Bases



Every ideal in $\mathbb{K}[\mathbf{a}, \mathbf{x}]$ has a finite set
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of generators $I = \langle g_1, \dots, g_N \rangle$
 $\{g_1, \dots, g_N\}$ is a **Gröbner basis** iff
$$g \in I \iff r_{g_1, \dots, g_N}(g) = 0$$

Gröbner Bases and Elimination

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And $I \cap \mathbb{K}[\mathbf{a}]$ is the (affine) ideal of the projection

Triangular sets



Triangular sets



- Characteristic set's method (J.F. Ritt - 1940's)

Triangular sets



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- Wu's method (Wenjun Wu - 1970's)

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Membership to I is easily tested

The “elimination” dictionary

Linear

Polynomial

The “elimination” dictionary

Linear

Gauss elimination

Polynomial

Gröbner Bases

The “elimination” dictionary

Linear

Gauss elimination

Triangulation

Polynomial

Gröbner Bases

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Triangulation

Determinants

Polynomial

Gröbner Bases

Triangular sets

Resultants

The “elimination” dictionary

Linear

Gauss elimination

Triangulation

Determinants

Cramer’s rule

...

Polynomial

Gröbner Bases

Triangular sets

Resultants

u-resultants

...

History Continues...

The rigorous foundation of the theory of algebraic varieties . . . can be formulated more simply than it has been done so far, without the help of **elimination theory**, on the sole basis of field theory and of the general theory of ideals in ring domains.

van der Waerden, 1926 First Paper on Algebraic Geometry



Fortunately....

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- MEGA, EACA, ACA, FoCM...

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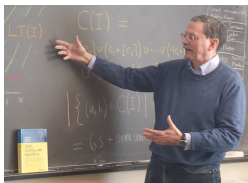
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Conclusion...

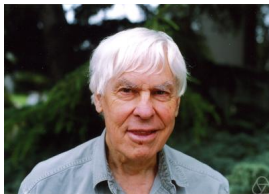
Elimination Theory is Important for
both **algorithmic** and **complexity**
aspects of polynomial system solving



It has led and still leads the
development of Algebraic Geometry
and Commutative Algebra in the last
60 years



It also impacts several other areas of
Mathematics (Numerical Analysis,
Linear Algebra, Complexity
Theory,...)



Main tools of (classic) Elimination

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■ Gröbner Bases

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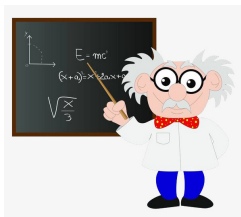
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Main tools of (classic) Elimination

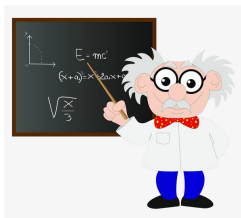
- Gröbner Bases
- Triangular Sets
- Resultants

$$\operatorname{Res}(f(x), g(x)) = \begin{vmatrix} a_n & a_{n-1} & \cdots & a_1 & a_0 & 0 & \cdots & 0 \\ 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & a_n & a_{n-1} & \cdots & a_1 & a_0 \\ b_n & b_{n-1} & \cdots & b_1 & b_0 & 0 & \cdots & 0 \\ 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & \cdots & 0 & b_n & b_{n-1} & \cdots & b_1 & b_0 \end{vmatrix}.$$

Syllabus

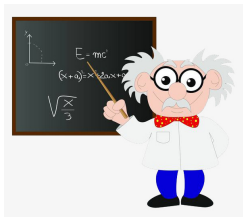


Syllabus



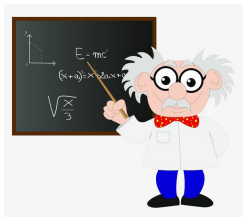
■ Univariate Resultant (afternoon)

Syllabus



- Univariate Resultant (afternoon)
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Syllabus



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- “Other” Resultants (Thursday)