

# Effective Pourchet's Theorem

Carlos D'Andrea

Nice 24/06/26



# Happy Birthday André!

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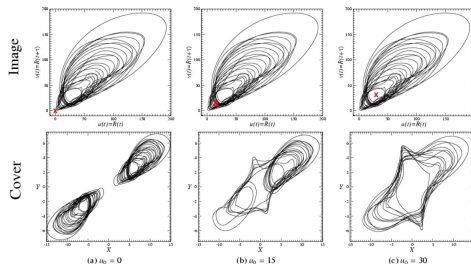


# Happy Birthday André!



# El roce hace el cariño

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Singularities of families of algebraic plane curves with “controlled” coefficients tend to the equidistribution  
 [Diaconis-Galligó]

JOURNAL ARTICLE

# QUANTITATIVE EQUIDISTRIBUTION FOR THE SOLUTIONS OF SYSTEMS OF SPARSE POLYNOMIAL EQUATIONS

Carlos D'Andrea, André Galligo and Martin Sombra

American Journal of Mathematics

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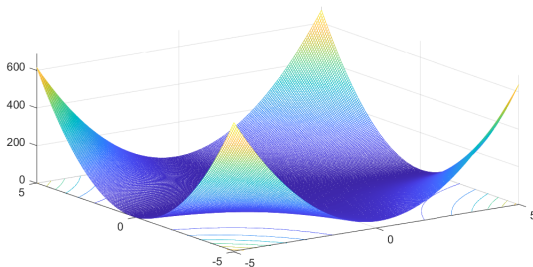
## Quantitative equidistribution of Galois orbits of small points in the $N$ -dimensional torus

Carlos D'Andrea, Marta Narváez-Clauss and Martín Sombra

We present a quantitative version of Bilu's theorem on the limit distribution of Galois orbits of sequences of points of small height in the  $N$ -dimensional algebraic torus. Our result gives, for a given point, an explicit bound for the discrepancy between its Galois orbit and the uniform distribution on the compact subtorus, in terms of the height and the generalized degree of the point.



# Today we have...



Yet another story involving  
polynomials

# Joint work with

- Ana Belén de Felipe – UPC
- Eulàlia Montoro – UB
- Joel Hurtado – UPC
- Teresa Cortadellas Benítez – URL

arXiv:2511.11783

# “The” general question

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Given a polynomial

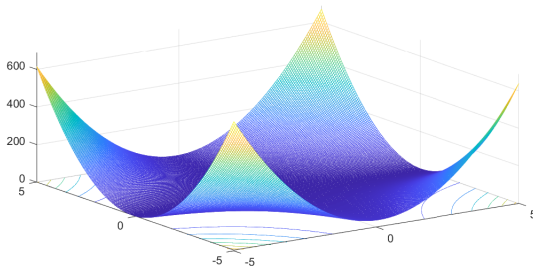
$$f(x_1, \dots, x_n) \in \mathbb{R}/\mathbb{Q}[x_1, \dots, x_n]$$

# “The” general question

Given a polynomial

$$f(x_1, \dots, x_n) \in \mathbb{R}/\mathbb{Q}[x_1, \dots, x_n]$$

How can we certify if  $f \geq 0$ ?



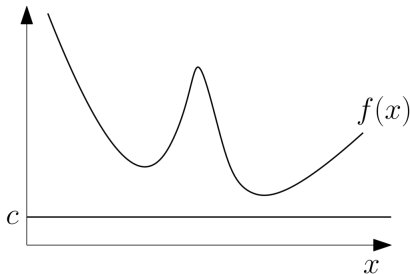
# In one variable...

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$$f(t) \geq 0 \quad \forall t \in \mathbb{R} \iff$$

# In one variable...

$$\begin{aligned} f(t) \geq 0 \quad \forall t \in \mathbb{R} &\iff \\ f(x) &= f_1(x)^2 + f_2(x)^2 \\ f_1(x), f_2(x) &\in \mathbb{R}[x] \end{aligned}$$





# Rational coefficients?



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$$f(t) \geq 0 \quad \forall t \in \mathbb{R} \iff$$

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$$f(t) \geq 0 \quad \forall t \in \mathbb{R} \iff$$

$$f(x) =$$

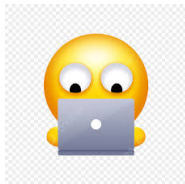
$$f_1(x)^2 + f_2(x)^2 + f_3(x)^2 + f_4(x)^2 + f_5(x)^2$$

Pourchet – 1971

# Is 5 optimal?

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$$x^2 + 7 = x^2 + 2^2 + 1^2 + 1^2 + 1^2$$



# Yves Pourchet

# Yves Pourchet



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ACTA ARITHMETICA  
XIX (1971)

Sur la représentation en somme de carrés des polynômes  
à une indéterminée sur un corps de nombres algébriques

par

Y. POURCHET (Paris)

Dans [7] Lazard démontre que tout polynôme défini positif à une indéterminée sur le corps des nombres rationnels est somme des carrés de huit polynômes sur ce corps. Je montre dans ce travail que ce nombre peut être ramené à cinq (corollaire 4 du théorème 2), sans d'ailleurs pouvoir être abaissé davantage (proposition 10).





# This talk – Effective Pourchet

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## ■ Input:

$$f(x) \in \mathbb{Q}[x], \quad f(t) \geq 0 \forall t \in \mathbb{R}$$

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## ■ Output: $f_1(x), \dots, f_5(x) \in \mathbb{Q}[x]$ ,

$$f(x) = f_1(x)^2 + \dots + f_5(x)^2$$



# Pourchet's original approach

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$$\begin{aligned} f(x) &= f_1(x)^2 + \dots + f_5(x)^2 \iff \\ f(x) &= f_{1p}(x)^2 + \dots + f_{5p}(x)^2 \\ \forall p &\in \{2, 3, 5, \dots\} \cup \{\infty\} \end{aligned}$$

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## ■ Local-Global Principle (Hasse-Minkowski)

# Pourchet's original approach

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- Local-Global Principle  
(Hasse-Minkowski)
- Non algorithmic



# An algorithm?

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$$p = \infty$$

Easy on  $\mathbb{R}$

$f(t) \geq 0 \iff y_1^2 + y_2^2 = f(x)$  has a  
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Proof:

$$x^2 + ax + b = (x - c)^2 + d^2 \text{ if } a^2 - 4b < 0$$

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$$(x - a)^{2k} = ((x - a)^k)^2 + 0^2$$

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$$(u^2 + v^2) \cdot (w^2 + z^2) = \alpha^2 + \beta^2$$

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**Any**  $f(x) \in \mathbb{Q}_p[x]$  is a sum of **at most four** squares if  $p \notin \{2, \infty\}$

# Almost all $p$

Any  $f(x) \in \mathbb{Q}_p[x]$  is a sum of **at most four** squares if  $p \notin \{2, \infty\}$   
Five squares are enough if  $p = 2$



# Sums of 4 squares are computable

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$$(x_1^2 + x_2^2 + x_3^2 + x_4^2)(y_1^2 + y_2^2 + y_3^2 + y_4^2) = (z_1^2 + z_2^2 + z_3^2 + z_4^2)$$

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Theorem (Pourchet, 71)

$$f(x) = f_1^2 + f_2^2 + f_3^2 + f_4^2 \text{ in } K[x] \iff$$

$$\blacksquare \text{lc}(f) = a_1^2 + a_2^2 + a_3^2 + a_4^2 \text{ in } K, \text{ and}$$

# Sums of 4 squares are computable

$$(x_1^2 + x_2^2 + x_3^2 + x_4^2)(y_1^2 + y_2^2 + y_3^2 + y_4^2) = (z_1^2 + z_2^2 + z_3^2 + z_4^2)$$

Theorem (Pourchet, 71)

$$f(x) = f_1^2 + f_2^2 + f_3^2 + f_4^2 \text{ in } K[x] \iff$$

- $\text{lc}(f) = a_1^2 + a_2^2 + a_3^2 + a_4^2$  in  $K$ , and
- $\forall p(x)$  prime divisor of  $f(x)$  with  $v_p(f)$  odd, there is a non trivial solution of  $x_1^2 + x_2^2 + x_3^2 + x_4^2 = 0$  in  $K[x]/(p(x))$

# Useful criteria

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$$x^2 + 7 = (x - \alpha) \cdot (x + \alpha) \text{ en } \mathbb{Q}_2[x]$$

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$$\begin{aligned} x^2 + 7 &= (x - \alpha) \cdot (x + \alpha) \text{ en } \mathbb{Q}_2[x] \\ &\implies \notin \sum_4 \mathbb{Q}[x]^2 \end{aligned}$$

# Useful criteria

$$x^2 + 7 = (x - \alpha) \cdot (x + \alpha) \text{ en } \mathbb{Q}_2[x] \\ \implies \notin \sum_4 \mathbb{Q}[x]^2$$

$$u \in \mathbb{Q}_2^2 \iff \\ u = 2^{2a}(8b + 1), \quad a \in \mathbb{Z}, \quad b \in \mathbb{Z}_2$$



# Algorithm

## **Pourchet's theorem in action: decomposing univariate nonnegative polynomials as sums of five squares**

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ISSAC 2023

# Sum of two squares

---

**Algorithm 1** Computing a decomposition of a polynomial as a sum of two squares

---

**Input:** A polynomial  $f \in \mathbb{Q}[x]$ , which is a priori known to be a sum of two squares in  $\mathbb{Q}[x]$ .

**Output:** Polynomials  $a, b \in \mathbb{Q}[x]$  such that  $a^2 + b^2 = f$ .

- 1: Construct the quadratic field extension  $\mathbb{Q}(i)/\mathbb{Q}$ .
- 2: Solve the norm equation

$$\text{lc}(f) = N_{\mathbb{Q}(i)/\mathbb{Q}}(x)$$

and denote a solution by  $a + bi \in \mathbb{Q}(i)$ .

- 3: Factor  $f$  into a product of monic irreducible polynomials

$$f = \text{lc}(f) \cdot p_1^{e_1} \cdots p_k^{e_k}.$$

- 4: **for** every factor  $p_j$ , such that the corresponding exponent  $e_j$  is odd **do**
- 5: Factor  $p_j$  over  $\mathbb{Q}(i)$  into a product  $p_j = g_j \cdot h_j$  with  $g_j, h_j \in \mathbb{Q}(i)[x]$ .
- 6: Set

$$a_j := \frac{1}{2} \cdot (g_j + h_j), \quad b_j := \frac{1}{2i} \cdot (g_j - h_j).$$

- 7: Update  $a$  and  $b$  setting:

$$a := aa_j + bb_j \quad \text{and} \quad b := ab_j - ba_j.$$

- 8: Update  $a$  and  $b$  setting:

$$a := a \cdot \prod_{j \leq k} p_j^{2\lfloor e_j/2 \rfloor} \quad \text{and} \quad b := b \cdot \prod_{j \leq k} p_j^{2\lfloor e_j/2 \rfloor}.$$

- 9: **return**  $a, b$ .

# Sum of three or four squares

---

**Algorithm 3** Initial solution: modular sum of squares

---

**Input:** An irreducible polynomial  $f \in \mathbb{Q}[x]$ , which is a priori known to be a sum of 3 or 4 squares.

**Output:** Polynomials  $h$  and  $g_1, \dots, g_4$  in  $\mathbb{Q}[x]$ , such that  $\deg h \leq \deg f - 2$  and  $fh = g_1^2 + \dots + g_4^2$ .

1: Construct the number fields:

$$K := \mathbb{Q}[x]/(f) \quad \text{and} \quad L := K(i).$$

2: Solve the norm equation

$$-1 = N_{L/K}(x)$$

and denote the solution by  $\xi = \bar{g}_1 + \bar{g}_2 i$ , where  $g_1, g_2 \in \mathbb{Q}[x]$  are polynomials of degree strictly less than  $\deg f$  and  $\bar{g}_j$  denotes the image of  $g_j$  under the canonical epimorphism  $\mathbb{Q}[x] \twoheadrightarrow K$ .

3: Set  $g_3 := 1, g_4 := 0$  and let  $h := (g_1^2 + \dots + g_4^2)/f$ .

4: **return**  $h, g_1, g_2, g_3, g_4$ .

---

# How to use this for 5 polynomials?



# How to use this for 5 polynomials?



$$f(t) > 0$$

# How to use this for 5 polynomials?



$$f(t) > 0 \dots f(t) - \left(\frac{1}{2^\ell}\right)^2 > 0$$

if  $\ell \gg 0$

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$$f(t) > 0 \dots f(t) - \left(\frac{1}{2^\ell}\right)^2 > 0$$

if  $\ell \gg 0$

$$f(x) - \left(\frac{1}{2^\ell}\right)^2 = f_1^2 + f_2^2 + f_3^2 + f_4^2??$$

# Algorithm 6



# Algorithm 6

---

**Algorithm 6** Reduction to a sum of 4 squares: odd valuation case

---

**Input:** A positive square-free polynomial  $f = c_0 + c_1x + \cdots + c_dx^d \in \mathbb{Q}[x]$ . The 2-adic valuations of the coefficients of  $f$  are  $k_j := \text{ord}_2 c_j$  for  $0 \leq j \leq d$ . Ensure  $k_d$  is odd. It is assumed that  $f$  is not a sum of 4 squares.

**Output:** A polynomial  $h \in \mathbb{Q}[x]$  such that  $f - h^2$  is a sum of 4 (or fewer) squares.

1: Find a positive number  $\varepsilon$  such that

$$\varepsilon < \inf \{f(x) \mid x \in \mathbb{R}\}.$$

2: Set  $l_1 := \lceil -1/2 \cdot \lg \varepsilon \rceil$ .

3: Set  $l_2 := \lceil -k_0/2 \rceil + 1$ .

4: Set

$$l_3 := \left\lceil \max \left\{ \frac{jk_d - dk_j}{2d - 2j} \mid 0 < j < d \right\} \right\rceil.$$

5: Initialize  $l := \max\{l_1, l_2, l_3\}$ .

6: **while**  $\gcd(d, 2l + k_d) \neq 1$  **do**

7:      $l := l + 1$ .

8: **return**  $h := 2^{-l}$ .

# Sum of 6 squares

---

**Algorithm 8** Decomposition of a nonnegative univariate rational polynomial into a sum of 6 squares

---

**Input:** A nonnegative polynomial  $f \in \mathbb{Q}[x]$ .

**Output:** Polynomials  $f_1, \dots, f_6 \in \mathbb{Q}[x]$  such that  $f_1^2 + \dots + f_6^2 = f$ .

- 1: **if**  $f$  is a square **then**
- 2:   **return**  $f_1 := \sqrt{f}, f_2 := \dots f_6 := 0$ .
- 3: **if**  $f$  is a sum of 2 squares {Use Observation 8 to check it} **then**
- 4:   Execute Algorithm 1 to obtain  $f_1, f_2 \in \mathbb{Q}[x]$  such that  $f_1^2 + f_2^2 = f$ .
- 5:   **return**  $f_1, f_2$  and  $f_3 := \dots f_6 := 0$ .
- 6: **if**  $f$  is a sum of 4 squares {Use [36, Theorem 17.2] to check it} **then**
- 7:   Execute Algorithm 5, to obtain  $f_1, \dots, f_4 \in \mathbb{Q}[x]$  such that  $f_1^2 + \dots + f_4^2 = f$ .
- 8:   **return**  $f_1, \dots, f_4$  and  $f_5 := f_6 := 0$
- 9: Compute the square-free decomposition of  $f = g \cdot h^2$ , where  $g, h \in \mathbb{Q}[x]$  and  $g$  is square-free.
- 10: Execute Algorithm 7 with  $g$  as an input to obtain  $g_1, g_2 \in \mathbb{Q}[x]$  such that  $g - g_1^2 - g_2^2$  is a sum of 4 squares in  $\mathbb{Q}[x]$ .
- 11: Execute Algorithm 5 to decompose  $g - g_1^2 - g_2^2$  into a sum of 4 squares in  $\mathbb{Q}[x]$ . Denote the output by  $g_3, \dots, g_6$ .
- 12: **return**  $f_1 := g_1 h, \dots, f_6 := g_6 h$ .

# Conjecture

---

**Algorithm 9** Reduction to a sum of 4 squares

---

**Input:** A positive square-free polynomial  $f = c_0 + c_1x + \dots + c_dx^d \in \mathbb{Q}[x]$ .

**Output:** A polynomial  $h \in \mathbb{Q}[x]$  such that  $f - h^2$  is a sum of 4 (or fewer) squares.

- 1: **if**  $f$  is a sum of 4 squares **then**
  - 2:     **return**  $h := 0$ .
  - 3: Set  $f_* := c_d + c_{d-1}x + \dots + c_0x^d$ .
  - 4: Find a positive number  $\varepsilon$  such that
$$\varepsilon < \inf\{f(x) \mid x \in \mathbb{R}\} \quad \text{and} \quad \varepsilon < \inf\{f_*(x) \mid x \in \mathbb{R}\}.$$
  - 5: Initialize  $l := \lceil -1/2 \cdot \lg \varepsilon \rceil$ .
  - 6: **while** True **do**
  - 7:     **if**  $f - 2^{-2l}$  is irreducible in  $\mathbb{Q}_2[x]$  **then**
  - 8:         **return**  $h := 2^{-l}$ .
  - 9:     **if**  $f - 2^{-2l}x^d$  is irreducible in  $\mathbb{Q}_2[x]$  **then**
  - 10:         **return**  $h := 2^{-l}x^{d/2}$ .
  - 11:      $l := l + 1$ .
-

# Our results

(CDDHM)

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- The conjecture works if  $\deg(f(x)) = 4k$

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- The conjecture works if  $\deg(f(x)) = 4k$
- It does not work for this family:

$$f_{k,N}(x) = \frac{4x^{2(2k+1)} + x^{2k+1} + 4}{N^2}$$

$$k = 0, \dots, N \in \mathbb{N} \text{ odd}, N > 64$$

# Extension of the method

(CDDHM)

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## Theorem

If  $f(x) \in \mathbb{Q}[x]$  of degree  
 $d = 2(2k + 1)$ ,  $k \in \mathbb{N}$ ,  $\ell \in \mathbb{N}$  such  
that  $f(t) - \frac{1}{2^{2\ell}}(t^2 + t + 1)^{2k}t^2 > 0 \forall t \in \mathbb{R}$



# Extension of the method

(CDDHM)

## Theorem

If  $f(x) \in \mathbb{Q}[x]$  of degree  
 $d = 2(2k + 1)$ ,  $k \in \mathbb{N}$ ,  $\ell \in \mathbb{N}$  such  
that  $f(t) - \frac{1}{2^{2\ell}}(t^2 + t + 1)^{2k}t^2 > 0 \forall t \in \mathbb{R}$   
then  $f(x) - \frac{1}{2^{2\ell}}(x^2 + x + 1)^{2k}x^2 \in \Sigma_4 \mathbb{Q}[x]^2$   
if  $f(0) \notin \mathbb{Q}_2^2$

# Missing case

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What if  $f(0) \in \mathbb{Q}_2^2$ ?

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$f(\alpha) \in \mathbb{Q}_2^2 \ \forall \alpha \in \mathbb{Q}_2 \Rightarrow f \in \mathbb{Q}[x]^2$ ?

# Missing case

What if  $f(0) \in \mathbb{Q}_2^2$ ?

$f(\alpha) \in \mathbb{Q}_2^2 \ \forall \alpha \in \mathbb{Q}_2 \Rightarrow f \in \mathbb{Q}[x]^2$ ?

$$4x^6 + 4x^3 + 9 = (1 + 2x^3)^2 + 8$$

# Missing case

What if  $f(0) \in \mathbb{Q}_2^2$ ?

$f(\alpha) \in \mathbb{Q}_2^2 \ \forall \alpha \in \mathbb{Q}_2 \Rightarrow f \in \mathbb{Q}[x]^2$ ?

$$4x^6 + 4x^3 + 9 = (1 + 2x^3)^2 + 8$$

## Conjecture

If  $f(x) \in \mathbb{Q}_2[x]$  is square-free and  $f(\mathbb{Q}_2^2) \subset \mathbb{Q}_2^2$  then it is a product of irreducibles of even degree in  $\mathbb{Q}_2[x]$

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- Semidefinite optimization ( $\mathbb{R}[x]$ )
- Over  $\mathbb{Q}[x]$

(Baldo-Krick-Mourrain 2025)

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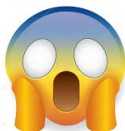
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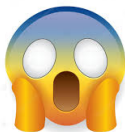


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## Hilbert's 17 Problem



# $\mathbb{R}$ vs $\mathbb{Q}$

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$$\begin{aligned} &40x_0^4 + 8x_0^2x_1^2 + 32x_0^2x_1x_2 + 64x_0^2x_1x_3 \\ &+ 16x_0^2x_2^2 + 16x_0^2x_2x_3 + 32x_0^2x_3^2 + 2x_1^4 \\ &+ 8x_1^2x_2^2 + 8x_1^2x_2x_3 + 16x_1x_2x_3^2 \\ &+ 8x_2^2x_3^2 + 8x_3^4 = f_1^2 + f_2^2 + f_3^2 + f_4^2 \end{aligned}$$



# $\mathbb{R}$ vs $\mathbb{Q}$

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over  $\mathbb{R}$

# $\mathbb{R}$ vs $\mathbb{Q}$

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over  $\mathbb{R}$  but cannot be written as any  
sum of squares over  $\mathbb{Q}$

# References

# References

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# Thanks!



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